

MATH 567 : Lecture 11 (02/13/2025)

Today: * unimodularity and total unimodularity

Recall: integral polyhedra...

Theorem 7 (Hoffman 1974) A rational polytope P is integral iff for all integral vectors \bar{w} , the optimal value of $\max \{\bar{w}^T \bar{x} \mid \bar{x} \in P\}$ is an integer. ↓
optimal $\bar{w}^T \bar{x}$

Proof (\Rightarrow) Let P be integral, i.e., all vertices are integral. Then $\max \{\bar{w}^T \bar{x} \mid \bar{x} \in P\}$ is integral for integral \bar{w} , as it occurs at a vertex.

(\Leftarrow) Let $\max \{\bar{w}^T \bar{x} \mid \bar{x} \in P\}$ be integral for all integral \bar{w} .

Let $\bar{v} = \begin{bmatrix} v_1 \\ v_2 \\ \vdots \\ v_n \end{bmatrix}$ be a vertex of P ,

and let \bar{w} be an integral vector such that \bar{v} is the unique optimal solution to $\max \{\bar{w}^T \bar{x} \mid \bar{x} \in P\}$.

We can assume $\bar{w}^T \bar{v} > \bar{w}^T \bar{u} + u_1 - v_1$ for all other vertices \bar{u} of P . We can scale \bar{w} by a large integer if needed (i.e., when $u_1 - v_1 < 0$).

Then we have $\bar{w}^T \bar{v} + v_1 > \bar{w}^T \bar{u} + u_1$ for all other vertices \bar{u} of P . This inequality gives the following result.

$\Rightarrow \bar{v}$ is the unique optimal solution for the objective function vector $\bar{w}' = [w_1+1, w_2, \dots, w_n]^T = \bar{w} + \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$, as

we get $\bar{w}'^T \bar{v} > \bar{w}'^T \bar{u}$ for all other vertices $\bar{u} \in P$.

By assumption, $\bar{w}'^T \bar{v}$ and $\bar{w}'^T \bar{u}$ are integral.

Also, $\bar{w}, \bar{w}' \in \mathbb{Z}^n \Rightarrow v_i \in \mathbb{Z}$.

We repeat the argument for v_2, v_3, \dots, v_n . □

Can extend result to unbounded pointed polyhedra easily, and also to polyhedra in general (i.e., not pointed).

While this theorem specifies an if and only if condition for P to be integral, it does not appear easy to check. We will have to certify integrality of the optimal value for all $\bar{w} \in \mathbb{Z}^n$, for which the optimum exists. But could we specify some easier to check conditions which guarantee integrality?

Unimodularity and Total Unimodularity (TU)

We assume A, \bar{b} are integral (in $A\bar{x} = \bar{b}$ or $A\bar{x} \leq \bar{b}$).

Def Let $A \in \mathbb{Z}^{m \times n}$ with full row rank ($\text{rank}(A) = m \leq n$).
 A is unimodular if each basis of A has determinant ± 1 .
 $\hookrightarrow B_{m \times m}$ submatrix of A with $\text{rank}(B) = m$.

$$A = [B \ N], B \in \mathbb{Z}^{m \times m}, \text{rank}(B) = m, \det(B) = \pm 1 \Rightarrow B^{-1} \in \mathbb{Z}^{m \times m}$$

For $\bar{b} \in \mathbb{Z}^m$, $A\bar{x} = \bar{b}$ has all integer solutions.

Recall: basic feasible solutions (bfs's) for $\max \{ \bar{c}^T \bar{x} \mid A\bar{x} = \bar{b}, \bar{x} \geq \bar{0} \}$
correspond to corner points (vertices) of $\{ \bar{x} \mid A\bar{x} = \bar{b}, \bar{x} \geq \bar{0} \}$.

$$A = [B \ N], \bar{x} = \begin{bmatrix} \bar{x}_B \\ \bar{x}_N \end{bmatrix}. \text{ Set } \bar{x}_N = \bar{0}, \text{ solve for } \bar{x}_B.$$

$$A\bar{x} = \bar{b} \Rightarrow B\bar{x}_B + N\bar{x}_N = \bar{b}.$$

$$\text{We get } \bar{x}_B = B^{-1}\bar{b}$$

$$\bar{x} = \begin{bmatrix} B^{-1}\bar{b} \\ \bar{0} \end{bmatrix} \text{ is a basic solution.}$$

If a basic solution \bar{x} satisfies $\bar{x} \geq \bar{0}$, then it's feasible,
and hence is a bfs. Each bfs corresponds to a vertex
(or corner point) of P .

We can use the idea of bfs to give a correspondence between unimodularity of A and integrality of P .

Theorem 8 $A \in \mathbb{Z}^{m \times n}$, $\text{rank}(A) = m$. Then $P = \{\bar{x} \mid A\bar{x} = \bar{b}, \bar{x} \geq 0\}$

is integral for all $\bar{b} \in \mathbb{Z}^m$ iff A is unimodular. ensures
 P is pointed

$$A = [B \ N], \quad \bar{x} = \begin{bmatrix} \bar{x}_B \\ \bar{x}_N \end{bmatrix}, \quad \text{set } \bar{x}_N = 0 \Rightarrow \bar{x}_B = \bar{B}' \bar{b} \in \mathbb{Z}^m.$$

If A is not unimodular, $\det(B) \neq \pm 1$ for some basis B of A , and hence the corresponding basic solution will not be integral for all $\bar{b} \in \mathbb{Z}^m$.

Note that $\bar{B}' \bar{b}$ might be integral for some $\bar{b} \in \mathbb{Z}^m$ in this case, e.g., when $|\det(B)| = 2$, and $\bar{b} \in 2\mathbb{Z}^m$. But the result will not hold for all $\bar{b} \in \mathbb{Z}^m$.

We now define a stronger (tighter) property of A , which guarantees integrality for polyhedra defined in more general forms. In particular, we would like to relax the requirement that $\text{rank}(A) = m$ (i.e., be able to consider more general "shapes" of A , e.g., more tall than wide, i.e., $m > n$).

Def A matrix A is **totally unimodular** (TU) if every square submatrix of A has determinant $-1, 0$, or 1 . In particular, $A_{ij} \in \{-1, 0, 1\} \forall i, j$.

e.g., $A = \begin{bmatrix} 3 & 2 \\ 1 & 1 \end{bmatrix}$ is unimodular, but not TU.

$B = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}$ is TU (and unimodular here).

$A = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{bmatrix}$ is not TU, as $\det(A_{1:3, 1:3}) = -2$.

Notice that we have to check all square submatrices in the worst case. But it turns out we could check whether a matrix is TU or not in polynomial time by Seymour's decomposition algorithm, which runs in $O(n^3)$ time. But no implementation is known. Another algorithm by Truemper runs in $O(n^5)$ time, but is implemented, though.

→ check out <https://discopt.github.io/cmr/> [Combinatorial Matrix Recognition (CMR) library]

Here is a result that connects total unimodularity and integrality of polyhedra.

Theorem 9 [Hoffman & Kruskal, 1956]: Let $A \in \mathbb{Z}^{m \times n}$. Then $P = \{ \bar{x} \mid A\bar{x} \leq \bar{b}, \bar{x} \geq \bar{0} \}$ is integral $\Leftrightarrow \bar{b} \in \mathbb{Z}^m$ for which $P \neq \emptyset$ iff A is totally unimodular (TU).

Proof We need a proposition first.

Proposition The following statements are equivalent.

- (i) A is TU.
- (ii) A^T is TU. \rightarrow determinant is preserved under transposes
- (iii) $[A \ I_m]$ is unimodular \rightarrow full row rank now!
- (iv) $\begin{bmatrix} A \\ -A \\ I_n \\ -I_n \end{bmatrix}$ is TU.

Back to proof of Theorem 9.

$P = \{ \bar{x} \mid A\bar{x} \leq \bar{b}, \bar{x} \geq \bar{0} \}$ is integral \Leftrightarrow
 $\text{Proj}_{\bar{x}}(P^z)$ is integral, where $P_z = \{ \bar{z} \mid [A \ I]\bar{z} = \bar{b}, \bar{z} \geq \bar{0} \}$
 is integral, where $\bar{z} = \begin{bmatrix} \bar{x} \\ \bar{s} \end{bmatrix}$. \rightarrow slack variables

Now apply Theorem 8, which gives that P_z is integral
 iff $[A \ I]$ is unimodular, which is true iff A is TU,
 by the proposition above. \square

Here is another characterization of TU and integral polyhedra.

Theorem 10 Let $A \in \mathbb{Z}^{m \times n}$. A is TU iff $P = \{\bar{x} \mid A\bar{x} \leq \bar{b}\}$ is integral $\nexists \bar{b} \in \mathbb{Z}^m$ for which $P \neq \emptyset$.

\bar{x} is uvs (unrestricted in sign) here. We can replace \bar{x} by $\bar{x}^+ - \bar{x}^-$, where $\bar{x}^+, \bar{x}^- \geq 0$, and use Theorem 9.

The upshot is that as long as the constraint matrix of the LP is TU, we're in good shape. The form of the LP—General or standard — does not matter.

Operations that preserve total unimodularity

1. Swap two rows (or columns).
 2. Taking transpose.
 3. Scaling a row/column by -1 .
 4. Pivoting, i.e., converting a column to a unit vector using EROs.
 5. Adding a zero row/column, or a singleton row/column with the single nonzero entry being ± 1 .
 6. Repeating a row/column.
- ⋮