

MATH 567 : Lecture 17 (03/06/2025)

- Today :
- * MIR cuts
 - * Knapsack cuts
 - * cover inequalities

Mixed-integer Gomory Cut (MIG cut)

Extension of MIR to higher dimensions. Let

$$X = \left\{ (\bar{x}, \bar{y}) \in \mathbb{R}_{\geq 0}^m \times \mathbb{Z}_{\geq 0}^n \mid \sum_{j \in N} a_j y_j + \sum_{j \in M} a_j x_j = \beta \right\}$$

where $N = \{1, 2, \dots, n\}$ and $M = \{1, 2, \dots, m\} + n$.
 $\Rightarrow N, M$ are index sets for \bar{y}, \bar{x} , resp.

IDEA: Derive a valid inequality of the form $y \leq x + \beta$ from the equation defining X , and apply MIR.

$$\sum_{j \in N} a_j y_j + \sum_{j \in M} a_j x_j = \beta$$

\nearrow we ignore terms with $a_j \geq 0$ ($x_j \geq 0 \forall j$)

$$\Rightarrow \sum_{j: a_j \leq \beta_L} \lfloor a_j \rfloor y_j + \sum_{j: a_j > \beta_L} a_j y_j + \sum_{a_j < 0} a_j x_j \leq \beta$$

\uparrow $\lfloor a_j \rfloor - a_j$

$$\Rightarrow \underbrace{\left(\sum_{j: a_j \leq \beta_L} \lfloor a_j \rfloor y_j + \sum_{j: a_j > \beta_L} \lfloor a_j \rfloor y_j \right)}_y \leq \beta + \underbrace{\left(\sum_{j: a_j > \beta_L} a_{ju} y_j - \sum_{a_j < 0} a_j x_j \right)}_x$$

We now apply MIR to get $y \leq \frac{1}{\beta_u} x + \lfloor \beta \rfloor$.

$$\Rightarrow \left(\sum_{a_j \leq \beta_L} \lfloor a_j \rfloor y_j + \sum_{a_j > \beta_L} \underbrace{\lfloor a_j \rfloor}_{\lfloor a_j \rfloor + 1} y_j \right) \leq \lfloor \beta \rfloor + \left(\sum_{a_j > \beta_L} \frac{a_j - \beta_L}{\beta_U - \beta_L} y_j - \sum_{a_j < 0} \frac{a_j}{\beta_U} x_j \right)$$

$$\Rightarrow \sum_{a_j \leq \beta_L} \lfloor a_j \rfloor y_j + \sum_{a_j > \beta_L} \left(\lfloor a_j \rfloor + \frac{\beta_U - a_j}{\beta_U} \right) y_j + \frac{1}{\beta_U} \sum_{a_j < 0} a_j x_j \leq \lfloor \beta \rfloor$$

is the MGR cut.

Note: If $|M|=0$, i.e., there are no x_j 's, the usual CG cut

$$\sum_{j \in N} \lfloor a_j \rfloor y_j \leq \lfloor \beta \rfloor.$$

But the MGR cut gives

$$\sum_{a_j \leq \beta_L} \lfloor a_j \rfloor y_j + \sum_{a_j > \beta_L} \left(\lfloor a_j \rfloor + \underbrace{\frac{\beta_U - a_j}{\beta_U}}_{\geq 0} \right) y_j \leq \lfloor \beta \rfloor,$$

which is stronger than the CG cut.

Wolsey (Integer Programming) calls the MIR cut as the "basic mixed integer inequality", and the special case of MGR cut with $|M|=1, |N|=2$ as the "MIR inequality".

Example

$$X = \left\{ (\bar{x}, \bar{y}) \in \mathbb{R}_{\geq 0}^2 \times \mathbb{Z}_{\geq 0}^3 \mid \underbrace{2x_1}_{a_4} - \underbrace{x_2}_{a_5} + \underbrace{\frac{10}{3}y_1}_{a_1} + \underbrace{y_2}_{a_2} + \underbrace{\frac{11}{4}y_3}_{a_3} = \underbrace{\frac{21}{2}}_{\beta} \right\}.$$

$$a_{1l} = \frac{1}{3}, a_{1u} = \frac{2}{3}, a_{3l} = \frac{3}{4}, a_{3u} = \frac{1}{4}, \beta_l = \beta_u = \frac{1}{2}.$$

$$\Rightarrow \frac{10}{3}y_1 + y_2 + \frac{11}{4}y_3 - x_2 \leq \frac{21}{2} \text{ is valid for } X.$$

Note that we have removed the $2x_1$ term from lhs.

$$\Rightarrow \left\lfloor \frac{10}{3} \right\rfloor y_1 + y_2 + \left(\left\lfloor \frac{11}{4} \right\rfloor + \frac{(\frac{1}{2} - \frac{1}{4})}{\frac{1}{2}} \right) y_3 - \frac{1}{2}x_2 \leq \left\lfloor \frac{21}{2} \right\rfloor$$

is valid for X .

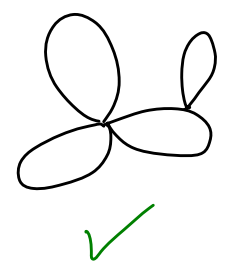
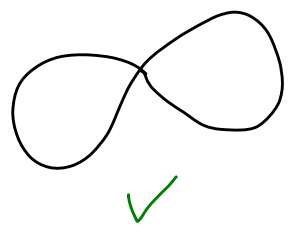
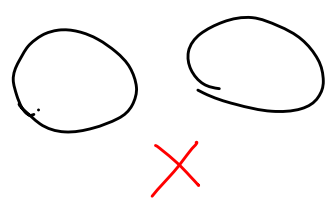
$$\Rightarrow 3y_1 + y_2 + \frac{5}{2}y_3 - 2x_2 \leq 10 \text{ is valid for } X.$$

Project 1: Hiker's tour problem (HTP)

$G = (V, E)$ directed graph

Find a circuit (closed walk) with following properties:

- * start and end at a given vertex;
- * do not have to visit every $v \in V$;
- * could visit a node more than once;
- * subtours are allowed as long as they are connected at vertices.



* $\sum_{(i,j) \in W} c_{ij} \geq L \leftarrow \text{data}$

Come up with formulations similar to the MTZ and subtour formulations for TSP.

Knapsack Cuts for pure 0-1 programs

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IDEA:

Given $\left\{ \begin{array}{l} \max \bar{c}^T \bar{x} \\ \text{s.t. } A\bar{x} \leq \bar{b} \\ \bar{x} \in \{0,1\}^n \end{array} \right\}$, pick $\bar{a}^T \bar{x} \leq \beta$ from $A\bar{x} \leq \bar{b}$,

generate cuts for $Y = \{ \bar{x} \in \{0,1\}^n \mid \bar{a}^T \bar{x} \leq \beta \}$, and add these cuts to the original IP.

Assume $a_i, \beta \in \mathbb{Z}$. WLOG, assume $a_i \geq 0 \forall i$ (in \bar{a}).

If $a_i < 0$, we could replace x_i with $(1-x_i)$ and a_i with $-a_i$ to get another inequality, for instance.

We define covers that capture the subsets of a_i that add to values larger than β (and hence "cover" it). If their sum is $> \beta$, we cannot have all the corresponding $x_j = 1$, which is the cut we are seeking.

Def $C \subseteq \{1,2,\dots,n\} = N$ is a **cover** if $\bar{a}(C) > \beta$, where $\bar{a}(C) = \sum_{i \in C} a_i$. Further, we say that C is a **minimal cover** if C is a cover, but $C \setminus \{i\}$ is not a cover $\forall i \in C$.

Example

Let $Y = \{ \bar{x} \in \{0,1\}^7 \mid \underline{11x_1 + 6x_2 + 6x_3 + 5x_4 + 5x_5 + 4x_6 + x_7 \leq 19} \}$



$C_1 = \{1, 4, 5\}$ is a minimal cover.

$C_2 = \{3, 4, 5, 6\}$ is a minimal cover.

$C_3 = \{3, 4, 5, 6, 7\}$ is a cover, but is not minimal.

Note that $a_3 + a_4 + a_5 + a_6 + a_7 = 21 > \beta = 19$. But so is $a_3 + a_4 + a_5 + a_6 (= 20)$.

Claim C is a cover $\implies \bar{x}(C) \leq |C| - 1$ is valid for Y .

Here, $\bar{x}(C) = \sum_{j \in C} x_j$.

$C_1: x_1 + x_4 + x_5 \leq 2$ is valid for Y . (1)

$C_2: x_3 + x_4 + x_5 + x_6 \leq 3$ is valid for Y (2)

$C_3: x_3 + x_4 + x_5 + x_6 + x_7 \leq 4$ is valid for Y , (3).

but (3) is weaker than (2), e.g., $x_3 + x_4 + x_5 + x_6 = 3.5$ satisfies (3), but violates (2).

Notice we added an extra variable (x_7) to the lhs of the \leq inequality with $x_7 \geq 0$, but also increased the rhs by 1. If we could add more nonnegative terms to the lhs while not changing the rhs, then we will strengthen the cut.

Def The extension of a cover C is

$$E(C) = \{j \notin C, j \in N \mid a_j \geq \max_{i \in C} \{a_i\}\} \cup C.$$

e.g., $E(\{3, 4, 5, 6\}) = \{1, 2, 3, 4, 5, 6\}.$

Claim $\bar{x}(E(C)) \leq |C| - 1$ is valid for Y .

So, $x_3 + x_4 + x_5 + x_6 \leq 3$ can be strengthened
(2)

to $x_1 + x_2 + x_3 + x_4 + x_5 + x_6 \leq 3$ (4).

Since $\bar{a}(C) > \beta$, and the added to C to obtain $E(C)$ are such that $a_j \geq \max_{i \in C} (a_i)$, the validity of the new cut follows.