

MATH 567: Lecture 21 (03/27/2025)

- * facial disjunctions
- * practical algorithm
- * rank of cuts
- * semidefinite relaxation

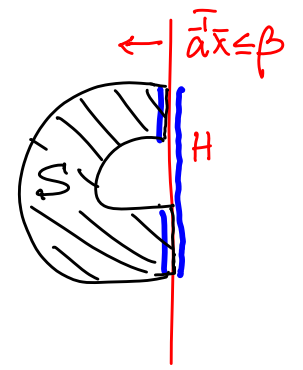
Def D_i is a **facial disjunction** w.r.t. K if D_{ij} are faces of K for $j=1, \dots, k_i$, i.e., $D_{ij} = \{\bar{x} \mid (\bar{a}^{ij})^T \bar{x} = \beta_{ij}\}$ where $(\bar{a}^{ij})^T \bar{x} \leq \beta$ is a supporting hyperplane of K .

Theorem 16 If D_1, \dots, D_p are facial disjunctions, then $K_p = K \cap_c (D_1 \cap \dots \cap D_p)$ in the theoretical algorithm.

Proposition 17 Let S be any set (possibly nonconvex), and $H = \{\bar{x} \mid \bar{a}^T \bar{x} = \beta\}$ is a hyperplane such that $\bar{a}^T \bar{x} \leq \beta \ \forall \bar{x} \in S$. Then $H \cap \text{conv}(S) = \text{conv}(H \cap S)$.

Proof (Theorem 16)

We show the result for $p=2, k_i=2, i=1,2$.



We need to show

$$(K \cap_c D_i) \cap_c D_j = K \cap_c (D_i \cap D_j)$$

$$\begin{aligned} (K \cap_c D_i) \cap_c D_j &= \text{conv}[\text{conv}(K \cap D_i) \cap D_j] \xrightarrow{D_{j_1} \cup D_{j_2}} \\ &= \text{conv}[(\text{conv}(K \cap D_i) \cap D_{j_1}) \cup (\text{conv}(K \cap D_i) \cap D_{j_2})] \\ &= \text{conv}[\text{conv}(K \cap D_i \cap D_{j_1}) \cup \text{conv}(K \cap D_i \cap D_{j_2})] \\ &\quad \text{by Proposition 17.} \end{aligned}$$

$$= \text{conv} [(K \cap D_i \cap D_{j_1}) \cup (K \cap D_i \cap D_{j_2})]$$

$$= \text{conv} [K \cap D_i \cap D_j] = K \cap_c (D_i \cap D_j). \quad \square$$

The theoretical algorithm might not work well in practice. Getting efficient descriptions of the convex hulls in each step might be difficult.

A practical Algorithm

Step 0: $K_0 = K = \{ \bar{x} \mid A\bar{x} \leq \bar{b} \}$

Step i : K_i : current relaxation, and
 ($i \geq 1$) \bar{x}^i : optimal solution over K_i

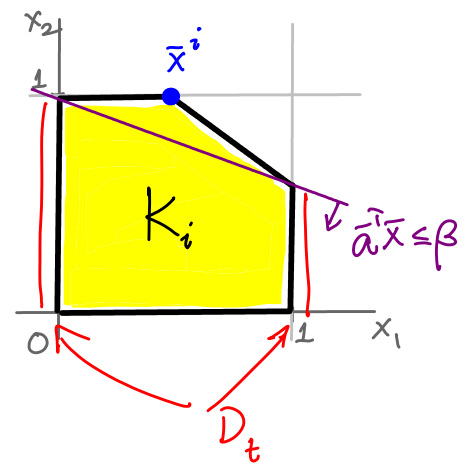
Find $\bar{a}^T \bar{x} \leq \beta$ such that

- ① $\bar{a}^T \bar{x} \leq \beta$ is valid for $K_i \cap D_t$ for some $t > i$.
- ② $\bar{a}^T \bar{x} \leq \beta$ is violated by \bar{x}^i (i.e., $\bar{a}^T \bar{x}^i > \beta$).

Set $K_{i+1} \leftarrow$ LP relaxation of $K_i \cap D_t \cap \{ \bar{x} \mid \bar{a}^T \bar{x} \leq \beta \}$.

How do we find (\bar{a}, β) ?

Solve an LP with \bar{a}, β as variables.



max $\beta - \bar{a}^T \bar{x}^i$ ↗ is given

s.t. $\left(\begin{array}{l} \bar{a}^T \bar{x} \leq \beta \text{ is derivable from } K_i \cap D_t \\ \text{by Farkas' lemma} \end{array} \right) \text{ --- (1)}$

normalization constraint : e.g.,

$\sum_{i=1}^n |a_{i1}| + |\beta| \leq 1.$ → we can linearize this constraint

Without the normalization, the LP could be unbounded. If (β, \bar{a}^T) works, then $100\beta - 100\bar{a}^T \bar{x}^i$ gives a bigger separation. And so does $10000\beta - 10000\bar{a}^T \bar{x}^i$.

For (1), we will use variables representing the multipliers for deriving the constraint $\bar{a}^T \bar{x}^i \leq \beta$ from $K_i \cap D_t$.

Q. How good is any such cutting plane algorithm? First, we define rank of cuts.

Rank of cuts

$K = \{ \bar{x} \mid A\bar{x} \leq \bar{b} \}$: All inequalities in $A\bar{x} \leq \bar{b}$, or derivable from $A\bar{x} \leq \bar{b}$ are **rank 0 cuts** (or inequalities).

Let $\bar{a}^T \bar{x} \leq \beta$ be valid for K . Then the CG cut $[\bar{a}^T] \bar{x} \leq [\beta]$ is a **rank 1 CG cut**.

Combining some (or all) rank-1 CG cuts and applying the CG procedure again gives me a **rank 2 CG cut**.

Similar notion of rank can be defined for the LS-procedure, MIG cuts, etc.

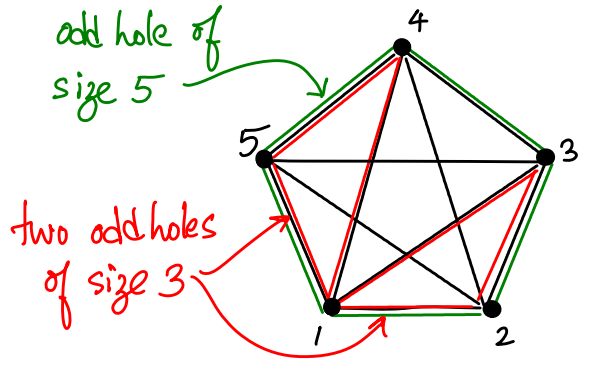
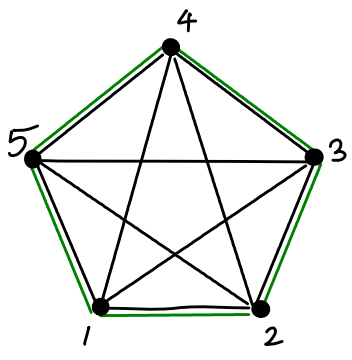
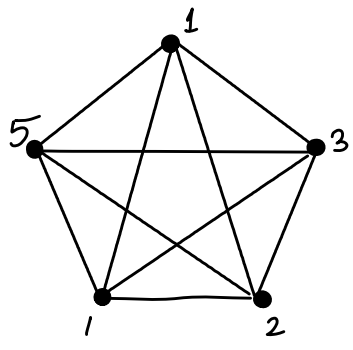
Ideally, we want to derive inequalities with small rank, so that we do not have to apply the cut-procedure too many times.

Back to LS Procedure

Q: How many steps to generate a good inequality?

Example: node packing on complete graph with 5 nodes.

$$K = \left\{ \begin{array}{l} x_i + x_j \leq 1 \quad \forall (i,j) \\ 0 \leq x_i \leq 1 \quad \forall i \end{array} \right\}$$



complete graph

Node packing problem:

A good inequality is $x_1 + x_2 + x_3 + x_4 + x_5 \leq 1$ ———— ⊗

We can pick at most one node.

Recall the definitions of $M_i^{NL}(K)$, $M_j(K)$, and $P_j(K)$. We have

$$M_i^{NL}(K) \text{ specified as } \left\{ \begin{array}{l} (A\bar{x} - \bar{b})x_i \leq 0 \\ (A\bar{x} - \bar{b})(1-x_i) \leq 0 \\ x_i(1-x_i) = 0 \end{array} \right\}$$

We get odd hole inequalities of size 3:

$$x_1 + x_2 + x_3 \leq 1$$

$$x_1 + x_2 + x_4 \leq 1$$

⋮

and $x_1 + x_2 + x_3 + x_4 + x_5 \leq 2$, odd-hole inequality of size 5.

We do not get \otimes .

In $P_2(P_1(K))$, we get $x_1 + x_2 + x_3 + x_4 \leq 1$ — (2)

To see (2) is valid for $P_2(P_1(K))$, we verify that (2) is valid for $P_1(K) \cap (x_2=0 \vee x_2=1)$.

$P_1(K) \cap (x_2=0)$ gives $x_1 + x_3 + x_4 \leq 1$, which is there in $P_1(K)$.

$P_1(K) \cap (x_2=1)$ gives $x_1 + x_3 + x_4 \leq 0$, which is also valid for $P_1(K)$, since we originally (in K itself) have $x_i + x_j \leq 1 \forall (i,j)$. Hence $x_2=1$ immediately forces $x_j=0 \forall j \neq 2$.

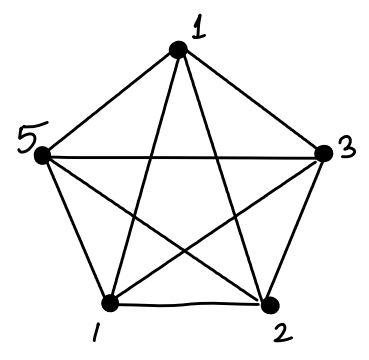
But we still do not get \otimes in $P_2(P_1(K))$.

$$P_3(P_2(P_1(K))) \text{ gives } x_1 + \dots + x_5 \leq 1 \text{ — } \otimes!$$

\Rightarrow LS-rank of \otimes is 3.

In general, LS rank of $x_1 + \dots + x_k \leq 1$ is $\leq k-2$, and is often $= k-2$.

Q. Could we derive $x_1 + x_2 + x_3 + x_4 + x_5 \leq 1$ — (*)
in one step? As a rank-1 cut.



YES!

Semidefinite Relaxation

Recall:

$$M_j^{NL}(K) \equiv \left\{ \begin{array}{l} (A\bar{x} - \bar{b})x_j \leq \bar{0} \\ (A\bar{x} - \bar{b})(1-x_j) \leq \bar{0} \\ x_j(1-x_j) = 0 \end{array} \right\}, j=1, \dots, n.$$

We linearize $M_j^{NL}(K)$ ($x_j^2 \leftarrow x_j, x_i x_j \leftarrow y_i$) to $M_j(K)$.

Define $M^{NL}(K) = \bigcap_{j=1}^n M_j^{NL}(K)$, and $M(K) = \bigcap_{j=1}^n M_j(K)$

↳ We could linearize all systems, and then take their intersection. Equivalently, we could take the intersection of the nonlinear systems, and then linearize.

Let $M_+(K) \equiv M(K) \cap \left\{ \bar{x} \mid \underline{(a_0 + \bar{a}^T \bar{x})^2 \geq 0} \forall [a_0, \bar{a}] \in \mathbb{R}^{n+1} \right\}$.

↳ $\begin{bmatrix} 1 & \bar{x}^T \\ \bar{x} & \bar{x}\bar{x}^T \end{bmatrix}$ is positive semidefinite

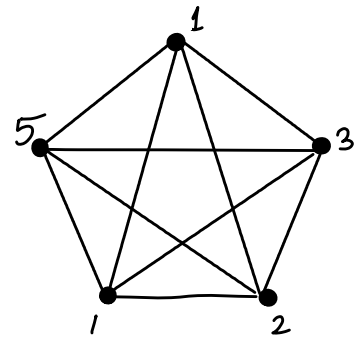
You would think this restriction is always satisfied! But imposing it explicitly makes the difference — see Notes to follow...

$M_+(K)$ is the semidefinite relaxation of the problem.

Def $A \in \mathbb{R}^{n \times n}$ is positive semidefinite (PSD) if $\bar{x}^T A \bar{x} \geq 0 \forall \bar{x} \in \mathbb{R}^n$. We write $A \succeq 0$

If A is PSD, all its eigenvalues are nonnegative.

Back to example on vertex packing



$$M_+(K) \text{ has } (1 - x_1 - x_2 - \dots - x_5)^2 \succeq 0$$
$$\Rightarrow 1 - 2 \sum_{i=1}^5 x_i + \sum_{i=1}^5 x_i^2 + 2 \sum_{i \neq j} x_i x_j \succeq 0$$

\downarrow
 x_i

\downarrow
 $= 0$

$$\Rightarrow 1 - \sum_{i=1}^5 x_i \geq 0, \text{ which is } (*)!$$

$$M(K) \text{ has } (x_i + x_j \leq 1) x_i \Rightarrow x_i x_j \leq 0 \text{ as } x_i^2 \leq x_i$$

$$M(K) \text{ also has } (-x_j \leq 0) x_i \Rightarrow x_i x_j \geq 0$$
$$\Rightarrow x_i x_j = 0.$$

Hence the semidefinite relaxation rank of $(*)$ is 1!

Notes

① $M^{NL}(K)$ contains inequalities for

$$X = \begin{pmatrix} 1 & \bar{x}^T \\ \bar{x} & \bar{x}\bar{x}^T \end{pmatrix} = \left[\begin{array}{c|cccc} 1 & x_1 & x_2 & \dots & x_n \\ \hline x_1 & x_1^2 & & & \\ x_2 & & x_2^2 & & x_i x_j \\ \vdots & & & \ddots & \\ x_n & & x_j x_i & & x_n^2 \end{array} \right].$$

Inequalities for X can be written as

$$B \bullet X \geq 0, \text{ where } A \bullet B = \text{trace}(A^T B)$$

② To get $M(K)$ from $M^{NL}(K)$, we replace x_i^2 by x_i and $x_i x_j$ by y_{ij} . Hence inequalities of $M(K)$ can be written as

$$C \bullet \begin{pmatrix} 1 & \bar{x}^T \\ \bar{x} & Y \end{pmatrix} \geq 0, \text{ where } \text{diag}(Y) = \bar{x}.$$

③ $(a_0 + \bar{a}^T \bar{x})^2 \geq 0$ can be written equivalently as

$$\begin{aligned} \begin{bmatrix} a_0 & \bar{a}^T \\ \bar{x} & \bar{x}\bar{x}^T \end{bmatrix} \begin{bmatrix} a_0 \\ \bar{a} \end{bmatrix} \geq 0 &\Rightarrow \begin{bmatrix} a_0 & \bar{a}^T \\ a_0 + \bar{a}^T \bar{x} \\ a_0 \bar{x} + (\bar{a}^T \bar{x}) \bar{a} \end{bmatrix} = \begin{bmatrix} a_0 & \bar{a}^T \\ (a_0 + \bar{a}^T \bar{x}) \\ (a_0 + \bar{a}^T \bar{x}) \bar{x} \end{bmatrix} \\ &= (a_0 + \bar{a}^T \bar{x})(a_0 + \bar{a}^T \bar{x}) \geq 0. \end{aligned}$$

In other words, $M_+(K)$ has $\begin{bmatrix} 1 & \bar{x}^T \\ \bar{x} & Y \end{bmatrix} \geq 0$ as added constraints.