

MATH 567 : Lecture 5 (01/23/2025)

Today: * representing sets
* representing functions

Proof of Theorem 1 (continued..)

Recall...

Theorem 1 \bar{x} satisfies $(*) \iff \exists (\bar{x}^1, \dots, \bar{x}^k, y_1, \dots, y_k)$ such that $(\bar{x}, \bar{x}^1, \dots, \bar{x}^k, y_1, \dots, y_k)$ satisfies $(*)$ -sharp.

$$\begin{aligned}
 & A_1 \bar{x}^1 \leq \bar{b}^1 y_1 \\
 & \vdots \\
 & A_k \bar{x}^k \leq \bar{b}^k y_k \\
 & \bar{x}^1 + \bar{x}^2 + \dots + \bar{x}^k = \bar{x} \\
 & y_1 + y_2 + \dots + y_k = 1 \\
 & y_i \in \{0, 1\}
 \end{aligned}$$

————— $(*)$ -sharp

Proof (\implies) : seen in the last lecture...

(\impliedby) WLOG, let $y_1=1, y_i=0, i=2, \dots, k$ in $(\bar{x}, \bar{x}^1, \dots, \bar{x}^k, y_1, \dots, y_k)$ that satisfies $(*)$ -sharp.

$$\begin{aligned}
 \implies & A_1 \bar{x}^1 \leq \bar{b}^1 \quad \text{and} \quad \bar{x}^1 + \dots + \bar{x}^k = \bar{x} \\
 & A_2 \bar{x}^2 \leq \bar{0} \\
 & \vdots \\
 & A_k \bar{x}^k \leq \bar{0}
 \end{aligned}$$

from Assumption 2

$$\implies A_1 \bar{x} = A_1 (\bar{x}^1 + \dots + \bar{x}^k) = A_1 \bar{x}^1 + A_1 \bar{x}^2 + \dots + A_1 \bar{x}^k \leq \bar{b}^1 + \bar{0} + \dots + \bar{0}$$

$$\implies A_1 \bar{x} \leq \bar{b}^1, \text{ i.e., } \bar{x} \text{ satisfies } (*).$$

□

Representing Sets in general

Q. In general, what all sets could we represent using 0-1 and/or general integer (G.I) variables?

We need a few new definitions to address this question. In particular, we will formally define a formulation - so far, we have been studying them informally as MIP models.

Def A set S is **bounded MIP-representable** (b-MIP-r) if

\exists matrices A, B, C, D and a vector \bar{f} such that

$$S = \{ (\bar{x}, \bar{y}) \in \mathbb{Z}^n \times \mathbb{R}^m \mid \exists (\bar{u}, \bar{v}) \in \mathbb{Z}^p \times \mathbb{R}^q \text{ such that } A\bar{x} + B\bar{y} + C\bar{u} + D\bar{v} \leq \bar{f} \}$$

$A\bar{x} + B\bar{y} + C\bar{u} + D\bar{v} \leq \bar{f}$ implies lower and upper bounds on \bar{x} and \bar{u} (the general integer (G.I) variables). → hence the "bounded" in b-MIP-r

The set $P = \{ (\bar{x}, \bar{y}, \bar{u}, \bar{v}) \in \mathbb{R}^{n+m+p+q} \mid A\bar{x} + B\bar{y} + C\bar{u} + D\bar{v} \leq \bar{f} \}$ is called a **formulation** of S .

Note that all variables are continuous in this formal definition of a formulation.

Def $P = \{ \bar{x} \in \mathbb{R}^n \mid A\bar{x} \leq \bar{b} \}$ is a polyhedron, where $A \in \mathbb{R}^{m \times n}$, $\bar{b} \in \mathbb{R}^m$, m, n are finite. → polyhedra are convex sets.
A bounded polyhedron is a **polytope**.

Examples of formulations

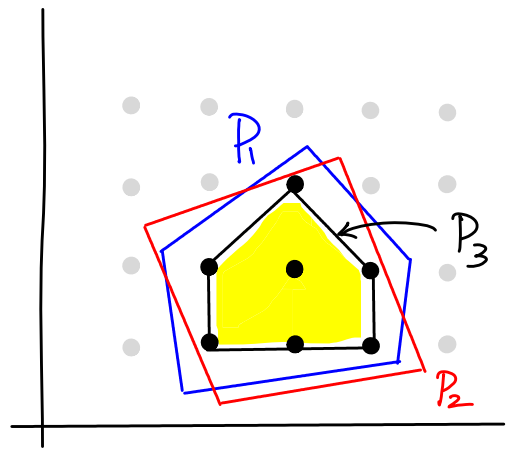
1. $S = \{ \bar{x} \in \{0,1\}^n \mid (x_1=1) \vee \dots \vee (x_n=1) \}$ has a formulation

$$P = \{ \bar{x} \in \mathbb{R}^n \mid 0 \leq x_i \leq 1 \forall i, x_1 + \dots + x_n \geq 1 \}$$

2. $S = \{ 7 \text{ lattice points shown as } \bullet \}$.

P_1, P_2, P_3 are all formulations for S .

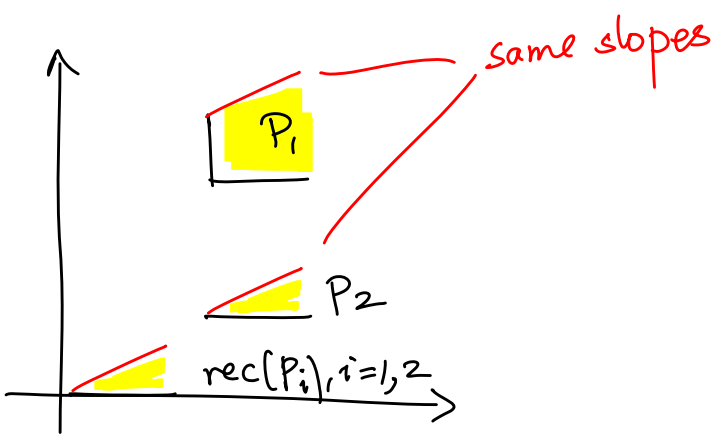
But P_3 is "better" than P_1 and P_2 . Note that P_3 is the convex hull of the lattice points that is S . If we want to maximize a linear function over S , we could do the same over P_3 instead. The same claim cannot be made for P_1 or P_2 . We will talk later about how to compare formulations.



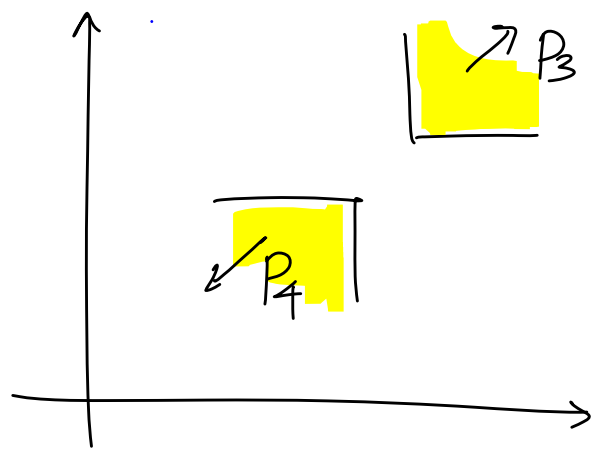
What kinds of sets S are b-MIP-r? Could we characterize them?

Theorem 2 (Jerashaw, Louve): S is b-MIP-r iff $S = P_1 \cup \dots \cup P_k$ for finite k , where P_i are polyhedra having the same recession cone i.e., $\text{rec}(P_i)$ is independent of $i, i=1, \dots, k.$

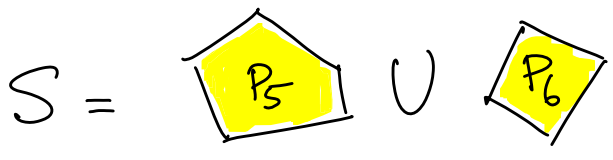
Here are some examples.



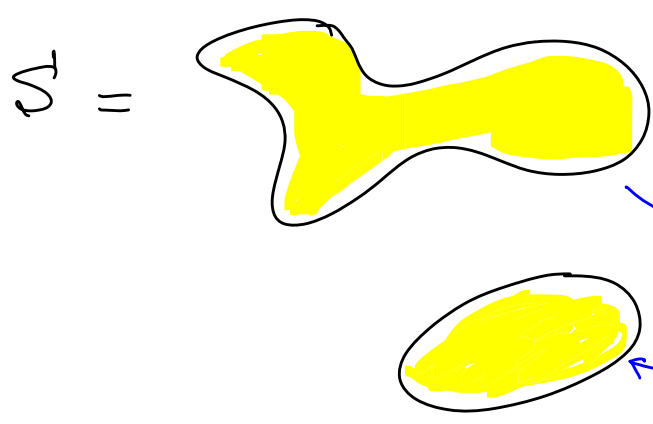
$P_1 \cup P_2$ is b-MIP-r.



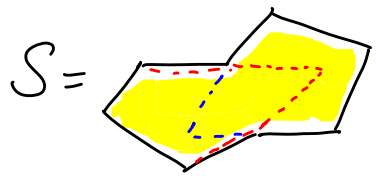
$P_3 \cup P_4$ is not okay as an MIP. (as $\text{rec}(P_3) \neq \text{rec}(P_4)$)



is okay as an MIP. ($\text{rec}(P_5) = \text{rec}(P_6) = \{ \vec{0} \}$)



is not okay as MIP, as it is not a polyhedron to start with. S being non convex is not crucial here — it could've been an ellipse, and the conclusion is the same.



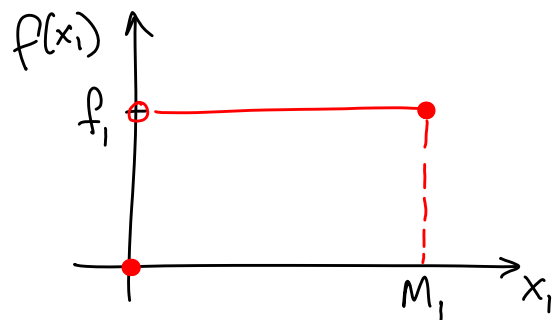
o.k. as MIP! \rightarrow union of two polytopes.

Can we use similar techniques to model functions, instead of sets?

Representing Sets v/s Functions

We already saw the case of fixed charge:

$$\left\{ \begin{array}{l} \min f(x_1) \\ \text{s.t. } 0 \leq x_1 \leq M_1 \\ A\bar{x} \leq \bar{b} \end{array} \right\} \text{ and we wrote an MIP for the same.}$$



It turns out we could use similar ideas to model some classes of functions appearing in certain optimization problems. Naturally, if $f(x_1)$ is nonlinear, e.g., x_1^3 or $\sqrt{x_1}$, we will not get an integer linear program! We need some definitions first.

Def Given $f: \mathbb{R}^n \rightarrow \mathbb{R}$, we define the **graph**, **epigraph**, and **hypograph** of f as follows.

$$\text{graph}(f) = \{ (z, \bar{x}) \in \mathbb{R}^{n+1} \mid z = f(\bar{x}) \},$$

$$\text{epi}(f) = \{ (z, \bar{x}) \in \mathbb{R}^{n+1} \mid z \geq f(\bar{x}) \}, \text{ and}$$

$$\text{hypo}(f) = \{ (z, \bar{x}) \in \mathbb{R}^{n+1} \mid z \leq f(\bar{x}) \}.$$

Notice that $\text{graph}(f)$, $\text{epi}(f)$, and $\text{hypo}(f)$ are sets, and we could consider when each of them is b-MIP-r — instead of talking about representability of $f(\cdot)$ itself.

Suppose we have $\left\{ \begin{array}{l} \min f(\bar{x}) \\ \text{s.t. } A\bar{x} \leq \bar{b} \end{array} \right\}$ where $\text{epi}(f)$ is b-MIP-r.

Then we can write $\left\{ \begin{array}{l} \min z \\ \text{s.t. } A\bar{x} \leq \bar{b} \\ (z, \bar{x}) \in \text{epi}(f) \end{array} \right\}$ as the MIP representation.

Since $\text{epi}(f)$ is b-MIP-r, we can write down an MIP representation of $\text{epi}(f)$, which completes the MIP model above.

Q. Why not require $\text{graph}(f)$ being b-MIP-r?

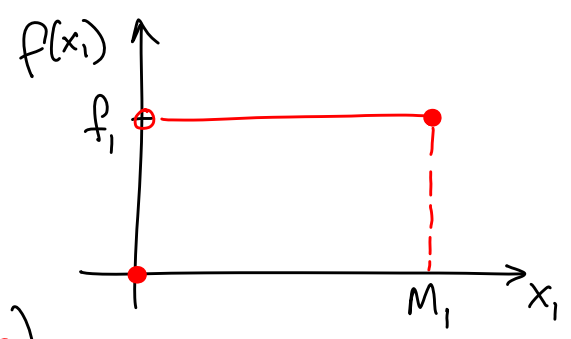
Theorem 3 $\text{graph}(f)$ is b-MIP-r iff both $\text{epi}(f)$ and $\text{hypo}(f)$ are b-MIP-r.

Example

The fixed charge function.

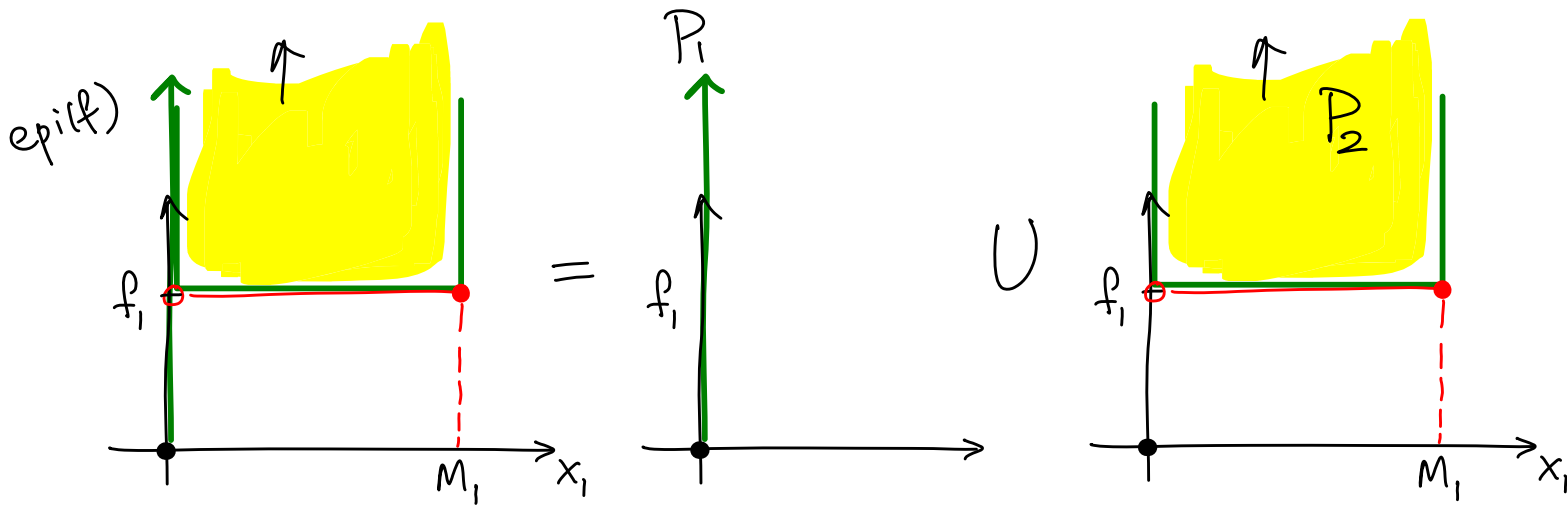
$\text{graph}(f)$, i.e., $f(x)$ as drawn, is not b-MIP-r.

$\text{graph}(f)$ is the union of origin (•) and the half-open line segment (○ — ●).



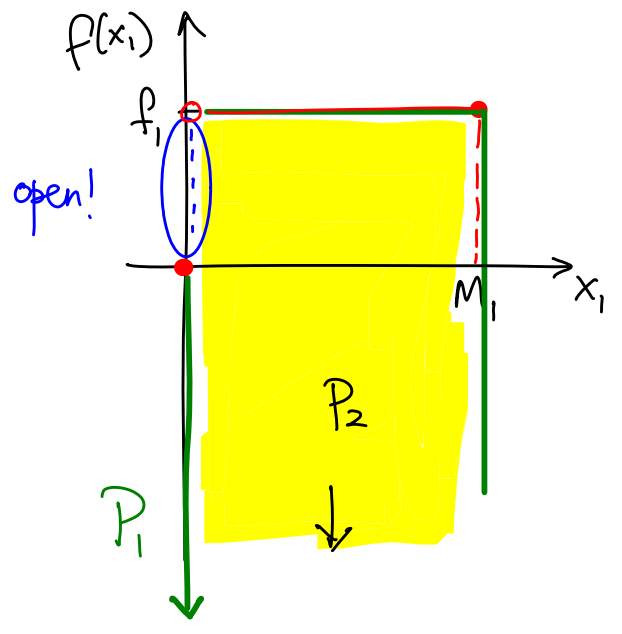
This second piece is not a polyhedron to start with.

But $\text{epi}(f)$ turns out to be b-MIP-r here, and $\text{hypo}(f)$ is not b-MIP-r at the same time.



Here $\text{rec}(P_1) = \text{rec}(P_2) = P_1$. Hence $\text{epi}(f)$ is b-MIP-r.

here, $\text{hypo}(f)$ is not b-MIP-r, as it is the union of P_1 and P_2 , where P_2 is not a polyhedron (same reason as that for $\text{graph}(f)$).

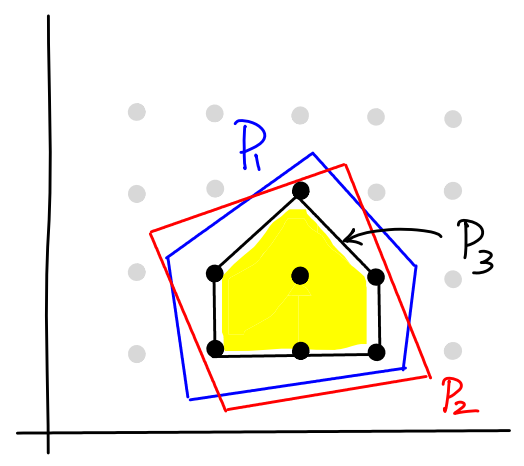


What if $f_1 < 0$ here? Then $\text{hypo}(f)$ will be b-MIP-r and $\text{epi}(f)$ will not be b-MIP-r (the roles are reversed).

Comparing Formulations

Q. What is a good/bad formulation?

We could say here P_3 is better than both P_1 and P_2 , but cannot compare P_1 to P_2 .
More generally, we want to compare formulations with different sets (and hence different numbers) of extra variables. To this end, we need to introduce some basic results.



→ pronounced "Farkash"; feasibility of alternative systems

Farkas' Lemma

① $\exists \bar{x} : A\bar{x} \leq \bar{b}$ then
 $A\bar{x} \leq \bar{b}$ implies $\bar{a}^T \bar{x} \leq \beta \iff$
 $\exists \bar{u} \geq 0$ such that $\bar{u}^T A = \bar{a}^T, \bar{u}^T \bar{b} \leq \beta$
 ↑
 multipliers using which we could derive $\bar{a}^T \bar{x} \leq \beta$ from $A\bar{x} \leq \bar{b}$

② $\exists \bar{x} : A\bar{x} \leq \bar{b} \iff \nexists \bar{u} \geq 0, \bar{u}^T A = \bar{0}^T, \bar{u}^T \bar{b} < 0.$
 (cannot derive $\bar{0}^T \bar{x} \leq -1$ from $A\bar{x} \leq \bar{b}$)